Internetworking With CISCO Routers & Switches



Course Outline





- Intro to Routing
- Router Interfaces
- Cisco Discovery Protocol (CDP)
- Routing
 - Static Routes
 - Distance Vector Routing
 - Link State Routing
 - Dynamic Routing
- Security Issues
- Advanced Topics

Graphic Symbols











Acces

server

ISDN switc h



switch

Network Multiswitch layer











Personal computer

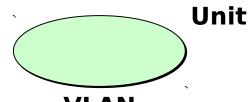


Data Service Unit/ Channel Service

Modem **Web Server**







VLAN (Color May Vary)



Hub



Network Cloud or Broadcast **Domain**

Ethern et

Fast Ethernet

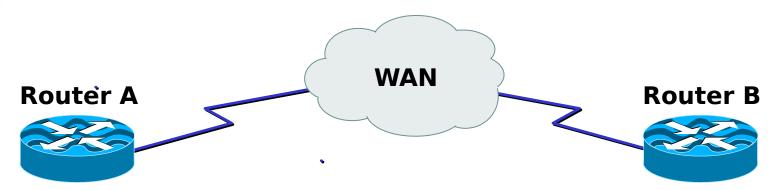
Serial Line

Circuit Switched Line

What do Routers do?





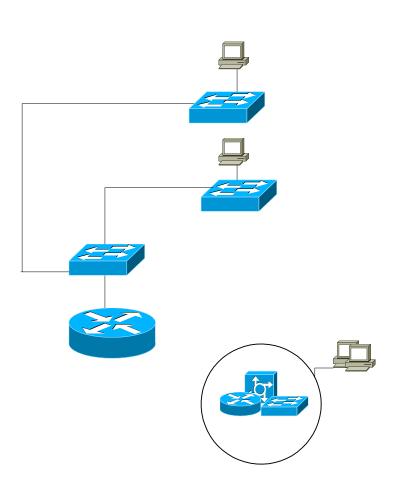


- Routers gather and maintain routing information to enable the transmission and reception of IP Datagrams
- Routing information is kept in a routing table
 - One entry for each known route
- Routers can create and maintain the routing table dynamically to accommodate network changes as they occur
- The rules for the exchange of routing information amongst routers are called routing protocols

What do Routers do?







- Broadcast Control
- Multicast protocol
- Optimal Path determination
- Traffic Management
- Connects to WAN services
- How do they do this?







	IFG	Minimum Valid Frame	Maximum Valid Frame	Bandwidt
Ethernet	96 bits	64 Bytes	1,518 Bytes	10 Mbps
Fast	96 bits	64 Bytes	1,518 Bytes	100 Mbps
Ethernet FDDI	0	34 Bytes	4,500 Bytes	100 Mbps
Token	4 bit	32 Bytes	16K Bytes	16 Mbps
Bixa	0	24 Bytes	1500 Bytes	128 Kbps
PRI	0	24 Bytes	1500 Bytes	1.472
T1	0	14 Bytes	4500 Bytes	Mbps 1.5 Mbps
АТМ	0	30 Bytes (AAL5)	16K Bytes (AAL5)	155 Mbps

LAN Interfaces: Broadcast and Multicast Traffic

- What's so bad about broadcasts?
 - Consume network bandwidth
 - Consume host station CPU capacity
- Sources of broadcast/multicast traffic
 - Clients looking for services
 - Apple Talk, Netware, Net BIOS, and TCP/IP clients
 - Servers announcing services
 - Routing protocol updates
 - Bridge Protocol Data Unit (BPDU) Frames
 - Forwarded by bridges and switches

LAN Interfaces:Broadcast and Multicast Traffic

- Controlling broadcast/multicast traffic
 - Rule of thumb:<20%
 broadcasts/multicasts per segment
 - Can be calculated with (# broadcasts)/(# packets input)
 - Limit maximum number of stations per segment
 - Guidelines

Ethernet: Performance Issues



- Measuring network utilization
 - Protocol analyzers
 - User complaints
 - Rule of thumb: shared Ethernet segments <40% utilization
- Improving network utilization
 - Segmenting with routers
 - Segmenting with switches

Routing





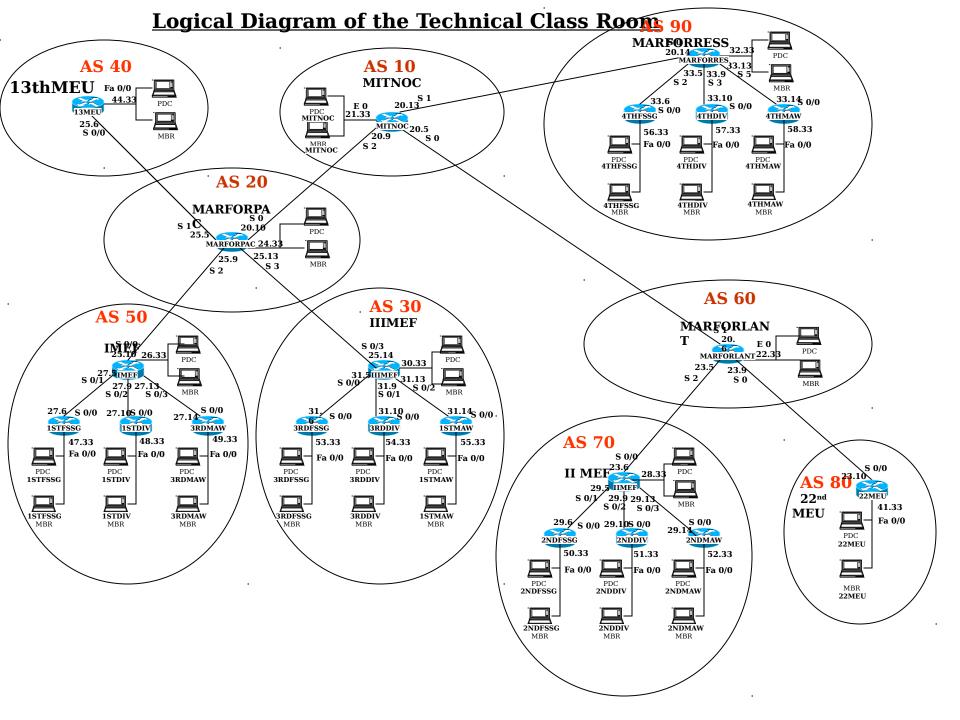
- Connected, Static, Default
- Distance Vector, Link-State
- Dynamic

IP Routing





- Routers learn routes by:
 - Directly connected networks
 - Routing information exchange with other routers
 - Static routes
- Default router
 - Every host should have a default router defined
 - ***Default router must be reachable



Routing Protocols



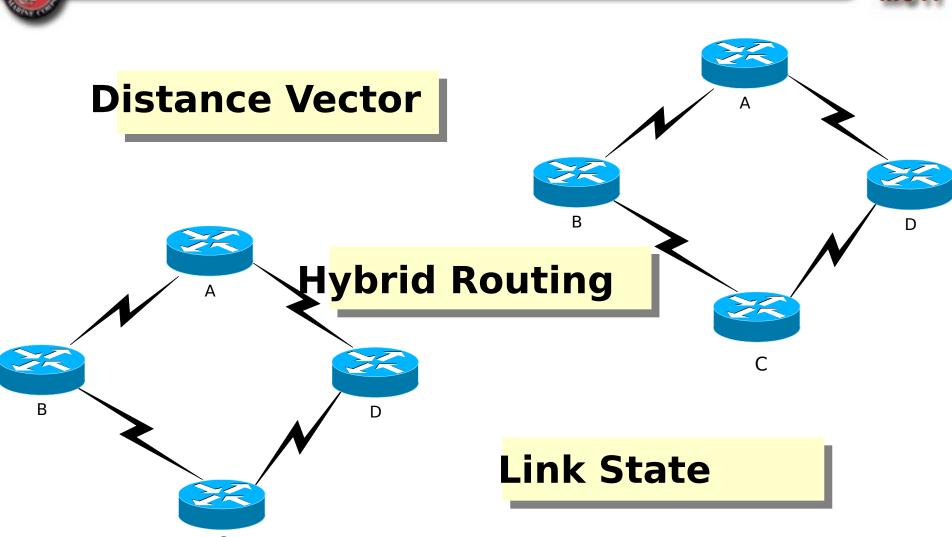


- Determine the "best" route to each destination network
- Distribute routing information amongst systems
- Distribute reach ability information amongst systems
- Interior routing protocol
 - Interior to an autonomous system
 - Under a common administration
 - Chosen by autonomous system's administrator
- Exterior routing protocol
 - Between autonomous systems
 - Not under a common administrator

Classes of Routing Protocols







Which Protocol?





Issue: Time to Convergence

Convergence occurs when all routers use a consistent perspective of network topology Change, routers must recompute routes, which disrupts routing

The process and time required for router reconvergence varies in routing protocols

Distance Vector vs. Link-State





Distance Vector	Link-State
 Views net topology from Adigh distance pecspectiven Foederto, rpetierdic updates: slow convergence Passes copies of routing table to neighbor routers 	 Gets common view of entire Getworker plotogy shortest path to Etherit rougeseed updates: faster Passesquices to the other routers

What is Best? It Depends





Issues	Concern	Example Questions
Technical	Performance to meet specific needs	Metrics adequate for network size? Any load sharing?
Business	Conformity with enterprise policies and priorities	Proven technology? Multi-vendor support? Standards based?
Operational	Simplicity of network setup and management	Easy to configure? Able to handle several routed protocols?

ch routing protocols, no one type fits all networks

Hybrid Routing



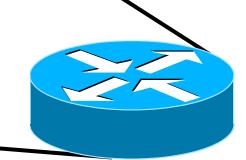


Choose a routing path based on distance vectors

Share attributes of both distance-vector and link-tate routing

alanced Hybrid Routing

Converge rapidly using changed-based updates



Autonomous System





- An autonomous system is a collection of networks under a common administration sharing a common routing strategy. An autonomous system may comprise of one or many networks, and each network may or may not have an internal structure (subnetting).
- The AS number, which is assigned by the NIC, is a 16-bit decimal number that is uniquely assigned.
- An assigned AS is required in order to run BGP, IGRP, or EIGRP.

The state of the s

Interior Routing Protocols



- EIGRP (Enhanced Interior Gateway Routing Protocol)
 - developed by Cisco
 - distance vector
 - sends out only updates when they happen vice whole tables every 30 seconds
 - supports VLSM
 - from global configuration
 HOSTNAME(conf)#router eigrp [AS]
 HOSTNAME(conf-router)#network [ip network]

Requirements for an Internet Routing Protocol





The protocol must be:

Scalabl

Stable

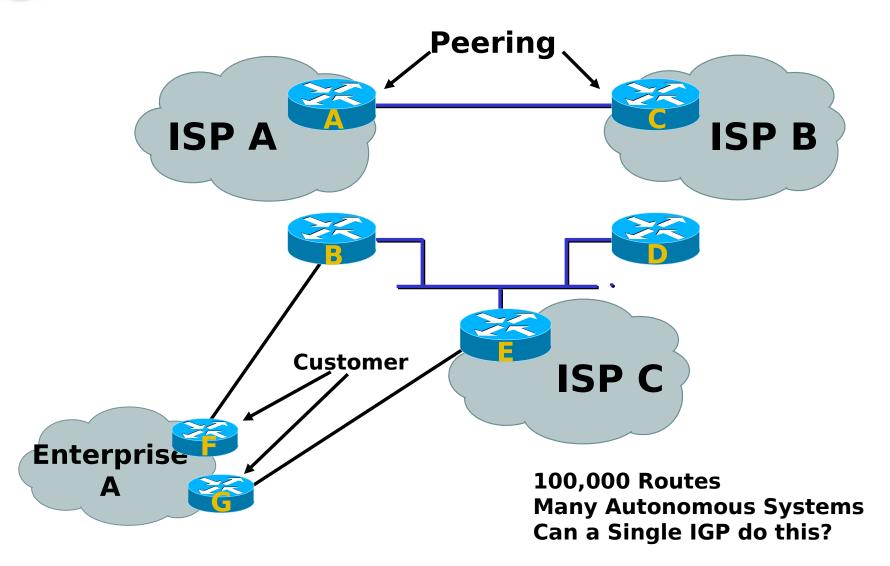
Flexible

...BGP is this Protocol

BGP Basics









Reasons for Using BGP



- 1: You need to scale your IGP
- 2: You're a multi-homed ISP customer and need to implement routing policy
- 3: You need to transit full Internet routes

Interior Vs. Exterior Routing





- Interior
 - Automatic discovery
 - Generally trust your IGP routers
 - Routes go to all IGP routers
- Exterior
 - Specifically configured peers
 - Connecting with outside networks
 - Set administrative boundaries



Why Do We Need an EGP?



- Scaling a large network—
 "Divide and Conquer"
 - Hierarchy
 - Periodic IGPs/Flooding
 - Isolate network stability
- Complex Policies
 - Control reachability to prefixes
 - Merge separate organizations
 - Connect multiple IGPs

Concept of Autonomous System



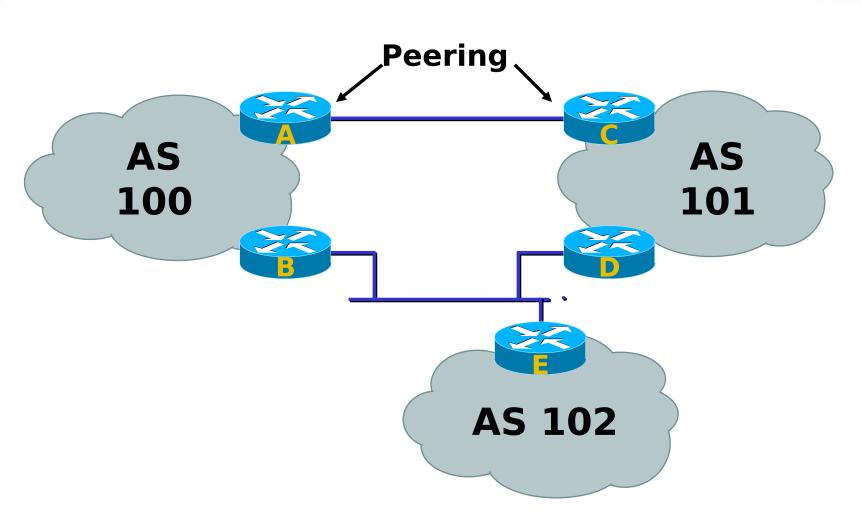


- A network(s) sharing the same routing policy
 - Possibly multiple IGPs
 - Usually under single administrative control
- Contiguous internal connectivity
- Numbering range form 1 to 65,535 globally unique—"AS Number"
 - Private range: 64512-65535



IGP of Each AS Is Hidden







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Stub Network



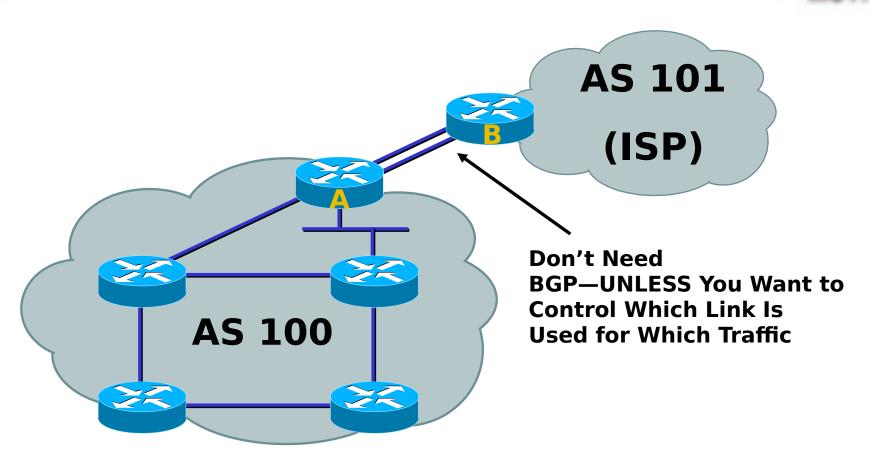


- No need for BGP
 - ISP advertises the stub network
 - Policy confined within ISP policy
- Default to the border

Stub Network







Multi-Homed Network



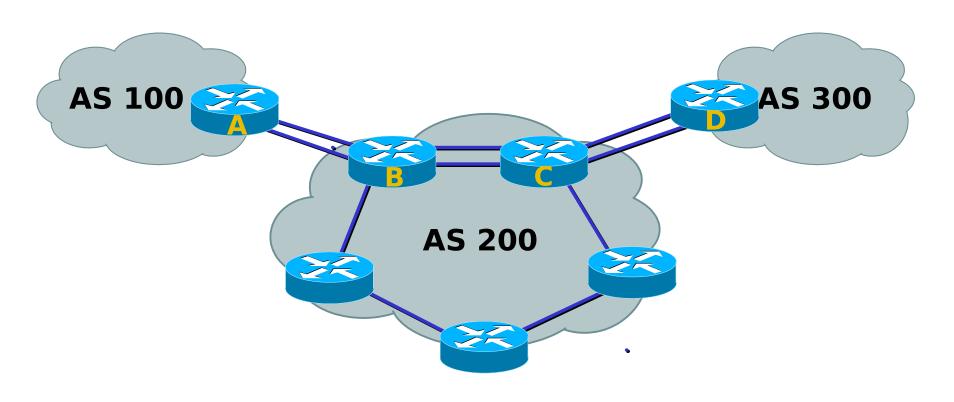


- Many situations possible
 - Multiple links to same ISP—without BGP
 - Secondary for only backup—without BGP
 - Loadshare between primary and secondary— without BGP
 - Selectively use different ISPs—need BGP

Multi-Homed Network





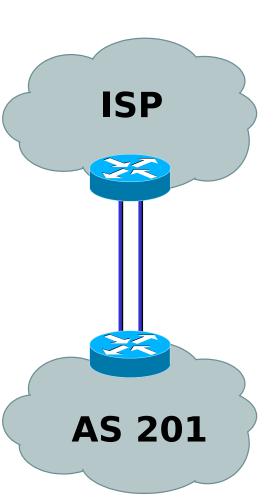


Can Still Use Default, UNLESS You Want to Selectively Use Either ISP for Optimal Performance

Multiple Links to the Same ISP I

MSTP

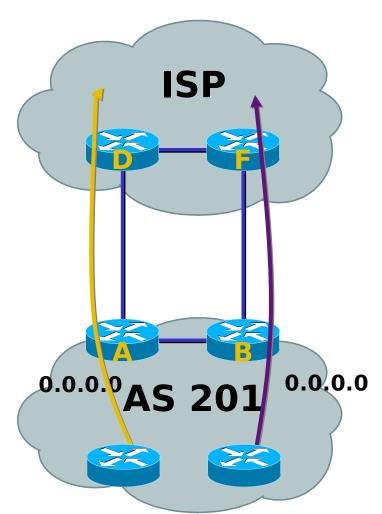
- Can still use default for outbound routing
- For inbound routing:
 - -Option1: ISP can use floating statics, or IGP to learn your routes and loadshare
 - -Option2: Can use BGP to loadshare



Multiple Links to the Same ISP



- Simplest scheme is to use two defaults
- Again, can use statics/IGP at borders, or use BGP

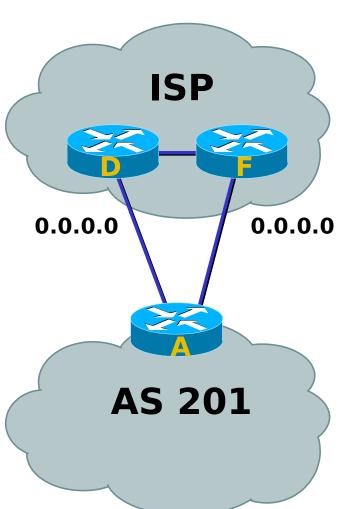


Multiple Links to the Same ISP

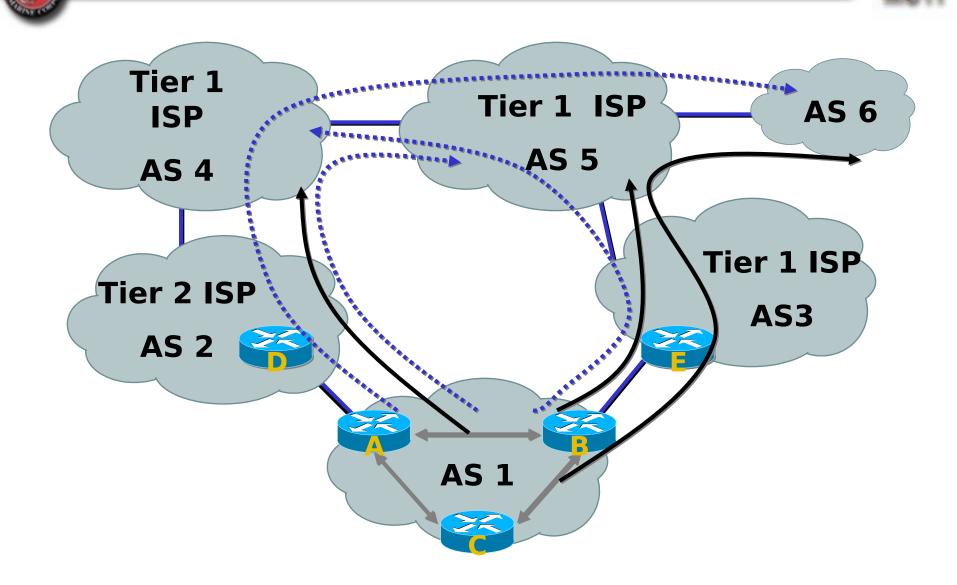




- Again, can just use two equal cost defaults to reach ISP
- Statics/IGP OR BGP to advertise your routes to ISP



Why Use BGP for Multi-Homing?



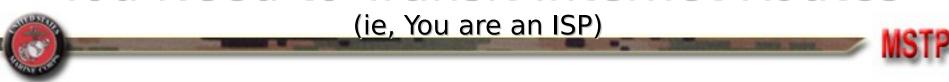


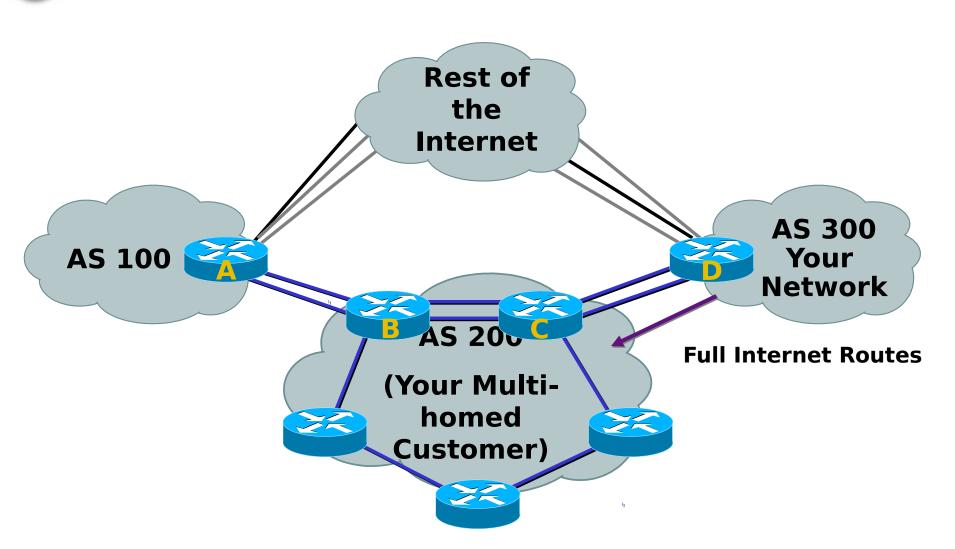
Reasons for Using BGP



- You need to scale your IGP
- You're a multi-homed ISP customer and need to implement routing policy
- You need to transit full Internet routes

You Need to Transit Internet Routes

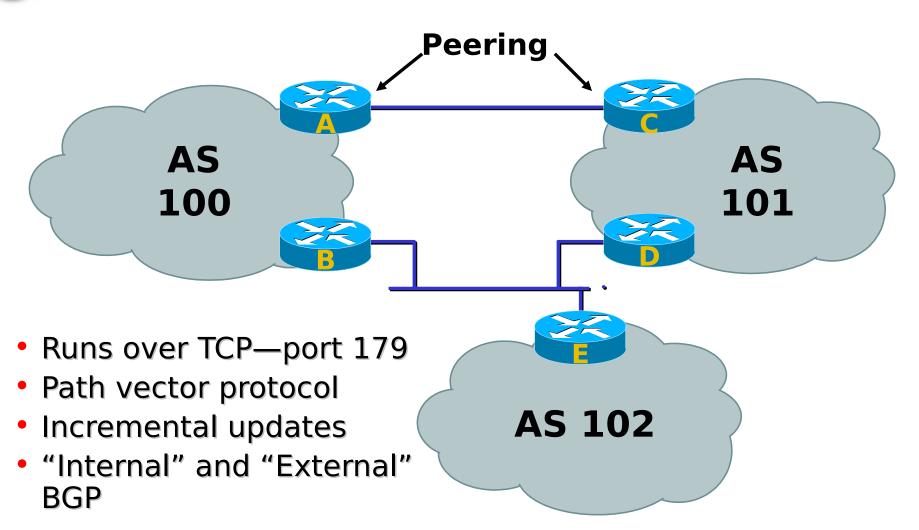




Back to Basics









General Operation



- Learns multiple paths via internal and external BGP speakers
- Picks THE bestpath, installs it in the IP forwarding table, forwards to EBGP neighbors (not IBGP)
- Policies applied by influencing the bestpath selection



Summary of Operation



- TCP connection established (port 179)
- Both peers attempt to connect—there is an algorithm to resolve "connection collisions"
- Exchange messages to open and confirm the connection parameters
- Initial exchange of entire table
- Incremental updates after initial exchange
- Keep alive messages exchanged when there no updates



Configuring Exterior Routing



- Configuring exterior routing protocols requires three sets of information.
 - A list of neighbor (or peer) routers with which to exchange routing information. This list is created with the neighbor router subcommand
 - A list of networks to advertise as directly reachable, created with the network router subcommand.
 - The AS number of the local router.
- Example BGP Configuration (> ver 4)
 - router bgp 110 (your local AS #)
 network 131.108.0.0 (network you wish to advertise)
 neighbor 131.108.100.2 remote-as 109 (IP of router in AS 109)

Default Routing





Manually configuring a static

```
( R1# (config) ip route 0.0.0.0 0.0.0.0 192.168.10.1
```

- Source a default route via IGP
 - Define a static route on one router
 - Redistribute static route into IGP

```
R1# (config) #ip route 0.0.0.0 0.0.0.0 192.168.10.1
R1# (config) #router rip
R1# (config-router) #redistribute static
```

- Specifying a default network
 - IGP will decide "best" route to default

 R1# (PONT) #ip default-network network-number

Redistributing Data





 Redistributing is the concept of passing unlike protocol information through different routing protocols. An example is IGRP will talk to all other IGRP clients, if you wish to also let those IGRP clients know about the manually added static routes on a router you would enter redistribute static as a router subcommand under the IGRP definition.

```
ip route 192.1.2.0 192.31.7.65 (static route)
ip route 193.62.5.24 255.255.255.248
192.31.7.65 (static route)
router igrp 110 (shares routing information with AS 110)
network 192.31.7.0 (locally connected network)
redistribute static (passes the two static routes as well)
redistribute rip (pass all RIP learned
```

Default Metric





The following example takes redistributed RIP metrics and translates them into EIGRP metrics with values as follows:

bandwidth = 1000 delay = 100 reliability = 250 loading = 100 mtu = 1500

router eigrp 109 network 131.108.0.0 redistribute rip default-metric 1000 100 250 100 1500

Routing Protocol Weights



The weight of a protocol helps the router to decide which is the

Directly Connected	0
Static	1
BGP (external)	20
EIGRP (internal)	90
IGRP	100
OSPF	110
RIP	120
EGP	140
EIGRP (external)	170
BGP (internal)	200
Unknown	255



Security Issues



- Password Security
- Keeping Out Unwanted Guests
- Allowing Certain Information

Passwords





- Enable
 - Enable password {password}
- Enable Secret
 - Enable secret {password}
- VTY
 - Line vty 0 4
 - Login
 - Password {password}
- Aux and Console
 - Line Con 0 (Aux 0)
 - Login
 - Password {password}

Password Security





- Password security is crucial for network integrity. It is recommended to mix the cases of the letters, and using "special characters" also helps in warding off hackers. Example: 2b,OR#2b.
- Never keep copies of configurations laying around that still display the passwords. Always edit the file after loading it onto a local drive.
- Outside administrators can still view most information without needing to know the Enable password. Always offer just a terminal login password unless they need to reconfigure as well.

Cisco Banners





- MOTD Message of the Day
 - Router#banner motd #Stop! Government site#
- Incoming
- Exec
- Login

Access Lists





- Access Lists determine what type of traffic is or is not allowed to travel through certain ports.
- Several Access Groups can be made to allow for different settings for particular ports or groups of ports.
- Some basic rules of access lists:
 - access commands are dependent on order of entry
 - if a packet meets any condition it is approved, else it is denied and is not sent out the port.
 - restrictions can be set to either incoming or outgoing traffic. Be sure to set the lists in the right direction.
 - It is usually best to use reverse logic for restrictions.
 That is to say, restrict all traffic, then set what is allowed.

Advance Topics



- Secondary Addresses
- Collapsing Backbones
- CIDR/Supernetting
- VLSM
- Queuing

Multiple Networks On The Same



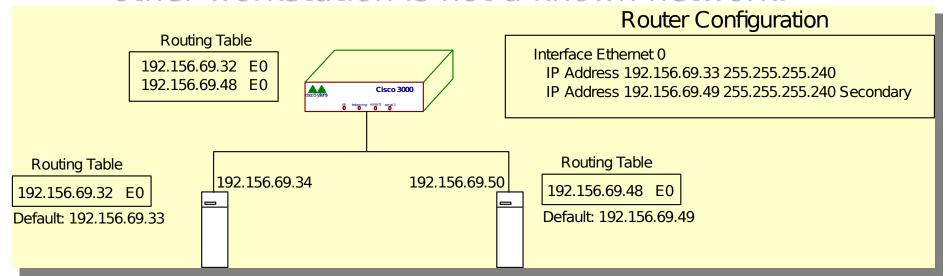
MSTP

- Should only be used when migrating
- Benefits
 - Allows for easy migrations
 - Works well for systems that use the broadcast address to advertise to clients, such as GCCS.
- Costs
 - Very inefficient at routing

Multiple Networks On The Same Line

MSTP

- Trace the network traffic between devices:
- Both workstations work fine when sending traffic to the router (default).
- When sending to each other, each workstation sends the traffic to the router first because the other workstation is not a known network.





Queuing



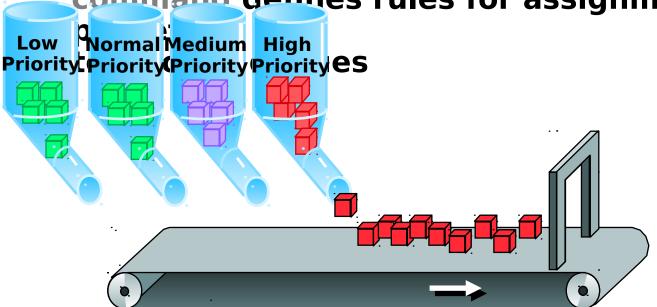
- The Cisco IOS implements four (now five) different queuing algorithms today:
 - First in, First Out (FIFO) Queuing
 - Priority Queuing
 - Custom Queuing
 - Weighted Fair Queuing
 - Interleave with Fragmentation
- Queuing occurs when network congestion occurs (i.e., the queue depth => 1), else all packets are sent as they arrive at the interface

Married Control

Priority Queuing



- Four queues: high, medium, normal and low
- Priority-list global and priority-group interface
 - command defines rules for assignment of

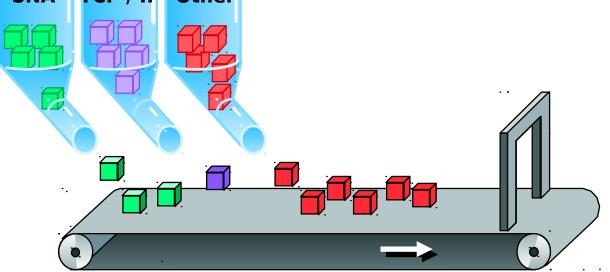


Custom Queuing





- Control % of interface bandwidth for specified traffic
- 17 output queues for each interface [16 configurable]
- · Queue-list global and custom-queue-list sinterface commands

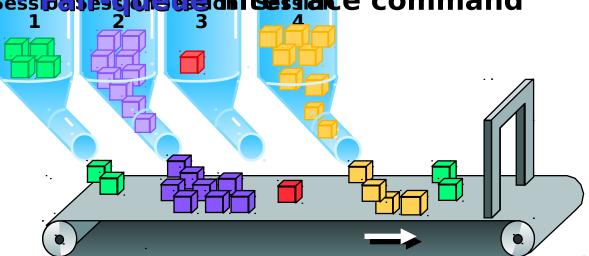


Weighted Fair Queuing





- Automatic traffic priority management
- Low-bandwidth sessions have priority over high-bandwidth sessions
- High-bandwidth sessions assigned weights
- Default for <2.048 Mbps interfaces



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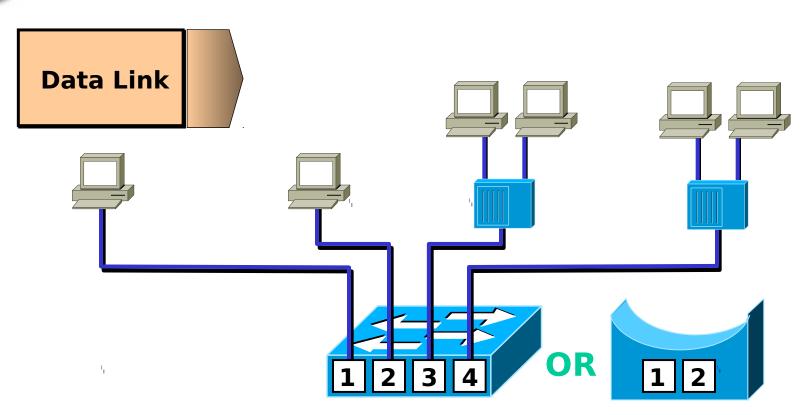
IOS Requirements



- Certain features may require specific IOS versions. i.e. QoS, IOS Telephony Service, etc
- Should be standard across the MEF.
 (Be aware of platform limitations!
 i.e. memory)

Switches, Bridges, &





- Each segment has its own collision domain
- All segments are in the same broadcast do

Ethernet: Collisions



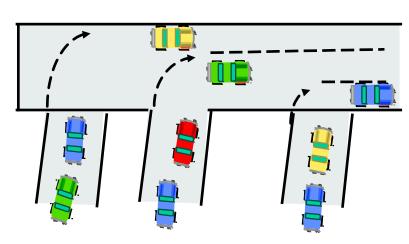


- A certain level of collisions are expected on CSMA/CD LANs
- Excessive collisions can result from faulty components or overloaded segments
 - Bad or excessively long cables
 - Bad NICs or transceivers
- Establishing a baseline is helpful to determine normal levels
- Collisions produce fragments that are <64 size of frame
- Local collisions
 - Occur on local LAN segment
 - Detected by circuitry in LAN interfaces
- Remote collisions

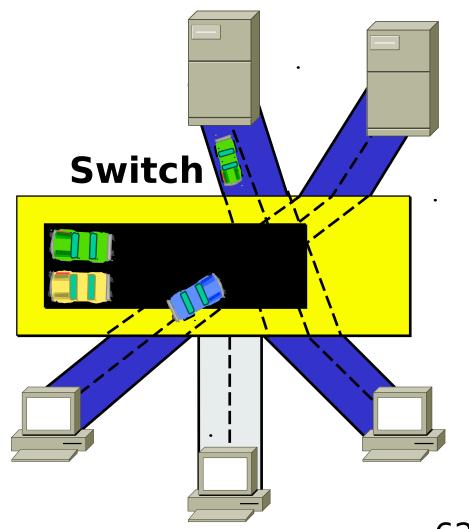
Advantage of Switches







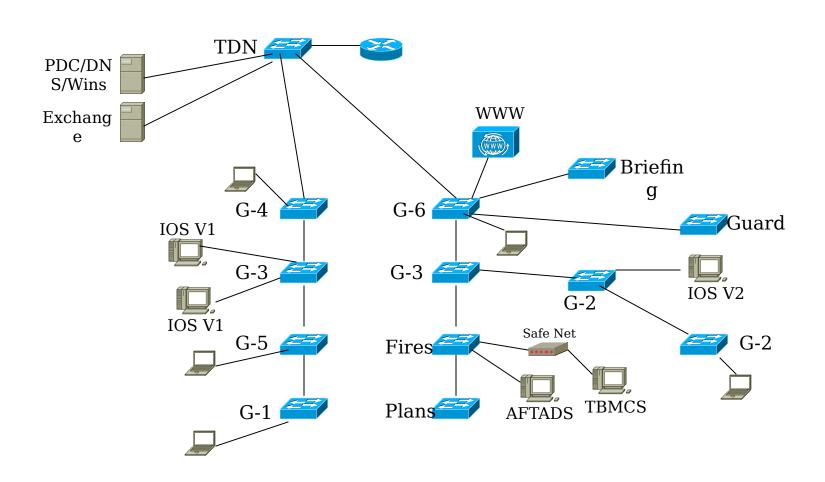
- Each segment has its own collision domain
- Broadcasts are forwarded to all segments





Switching Requirement

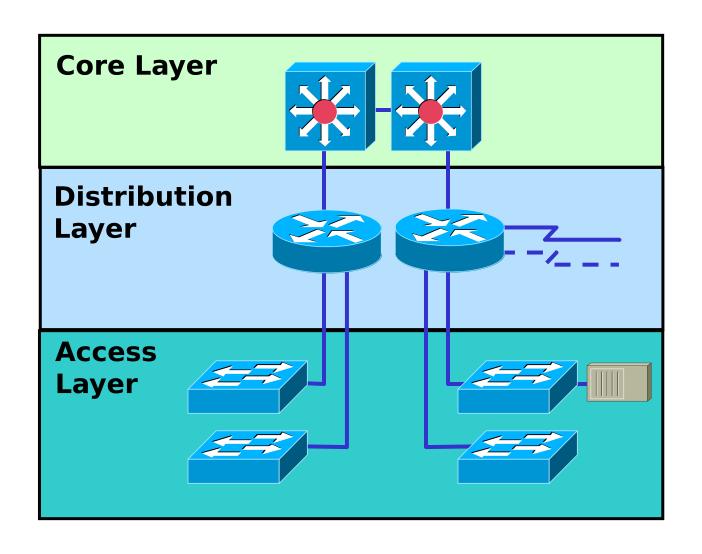




Cisco's Network Hierarchy



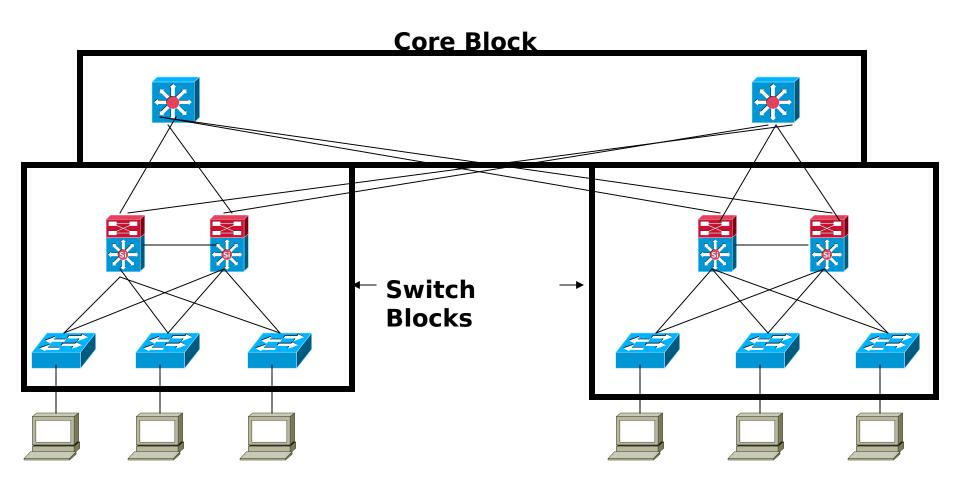




Building Block Method



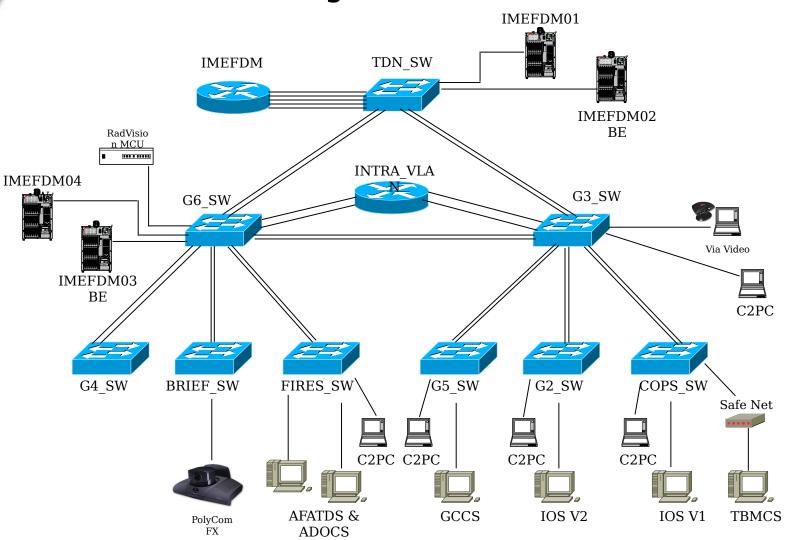




Switching Design (C2 Systems)







Switch Functions



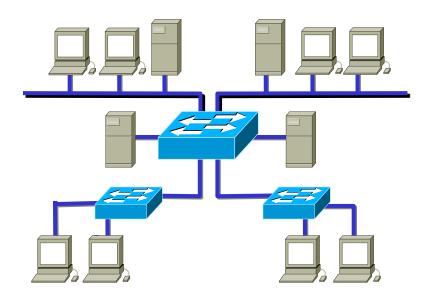


- Break Up Collision Domains
 - Layer 2 Switching is Hardware Based
 - Application-Specific Integrated Circuits (ASICs)
 - No modification to Layer 2 Header
- Provide Segmentation
 - Each Port is a segment
 - Can achieve Gigabit Speeds

Three Main Tasks





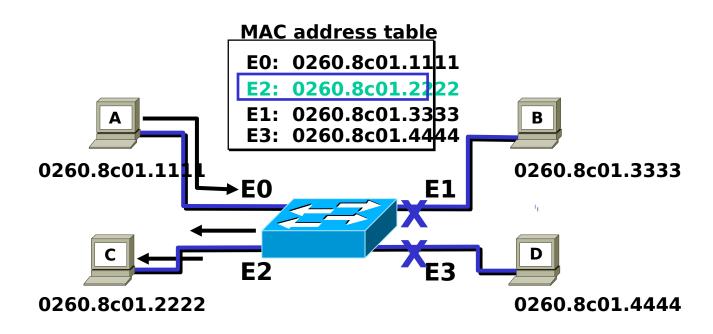


- Address learning
- Forward/filter decision
- Loop avoidance



Filtering Frames



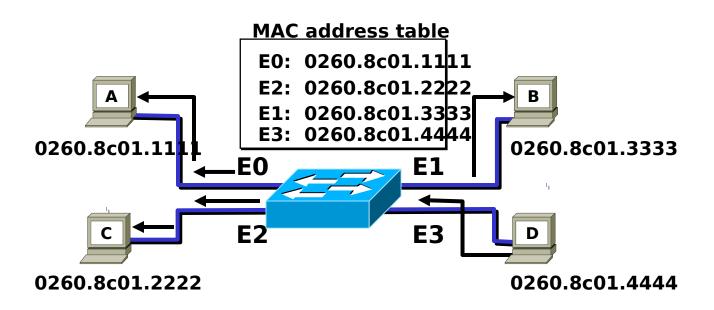


Station A sends a frame to station C Destination is known, frame is not flooded

Broadcast and Multicast





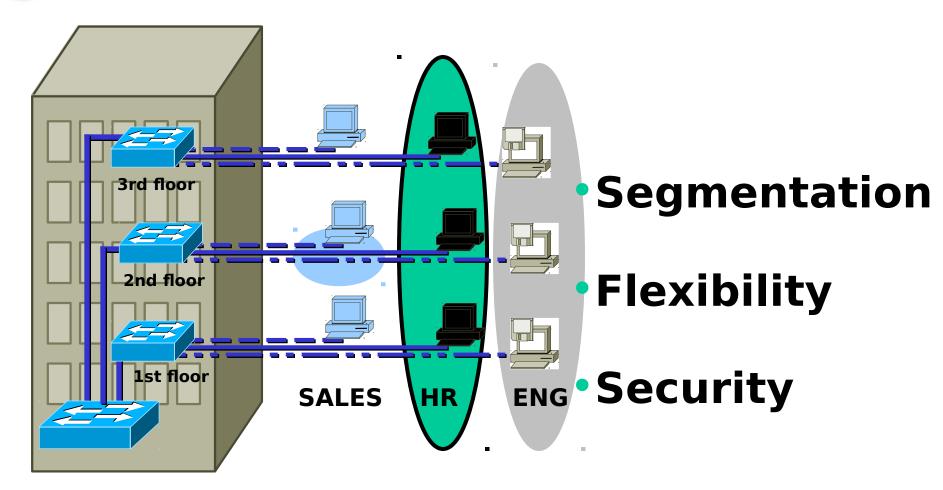


- Station D sends a broadcast or multicast frame
- Broadcast and multicast frames are flooded to all ports other than the

VLAN Overview



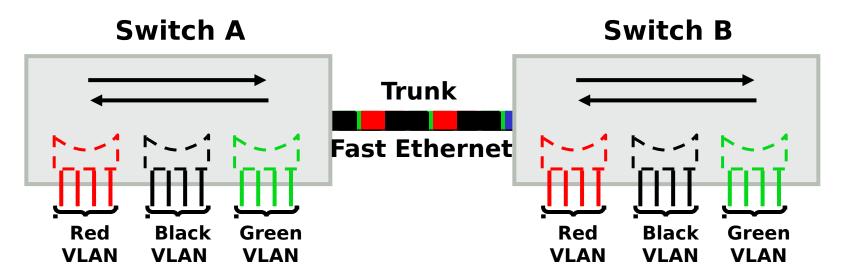




VLAN = A broadcast domain = Logical network (subn







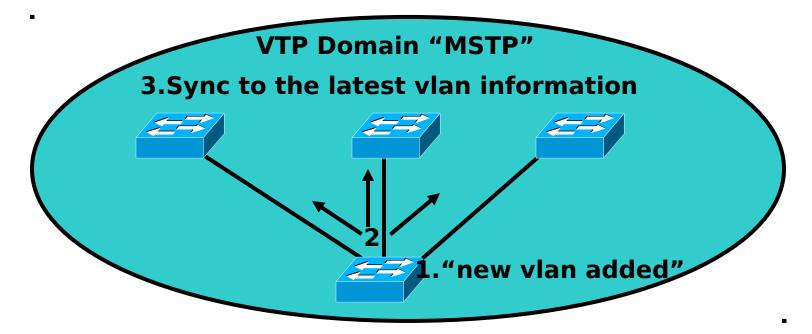
- Each logical VLAN is like a separate physical bridge
- VLANs can span across multiple switches
- Trunks carries traffic for multiple VLANs

VLAN Trunking Protocol (VTP)





- A messaging system that advertises VLAN configuration information
- Maintains VLAN configuration consistency throughout a common administrative domain
- VTP sends advertisements on trunk ports only
- Support mixed media trunks (Fast Ethernet, FDDI, ATM)



Client Mode

Synchronize

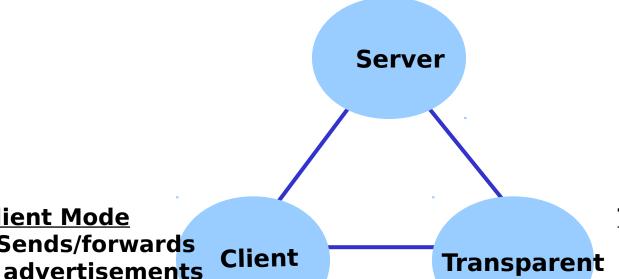
Not saved in

NVRAM

Sends/forwards

VTP Modes





Server Mode

- Create vlans
- Modify vlans
- Delete vlans
- Sends/forwards advertisements
- Synchronize **Domain**
- Saved in NVRAM

Transparent Mode

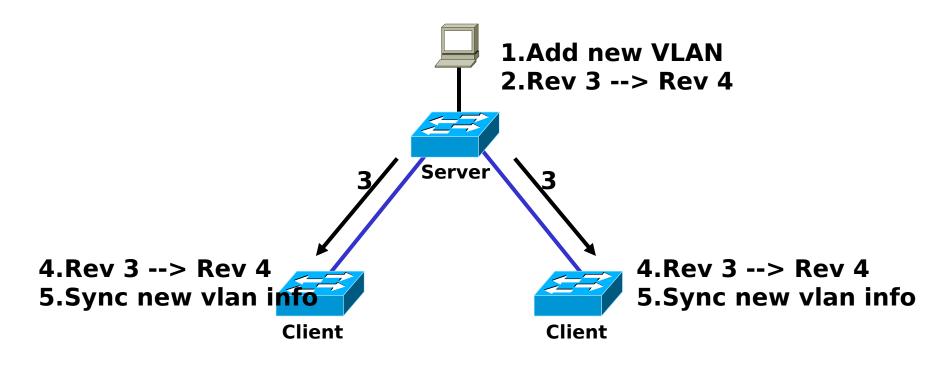
- Create vlans
- Modify vlans
- Delete vlans
- Forwards advertisements
- Does not synchronize
- Saved in **NVRAM**

How VTP Works





VTP advertisements are sent as multicast frames
VTP servers and clients synchronized to latest revision number
VTP advertisement are sent every five minutes or when there is a change

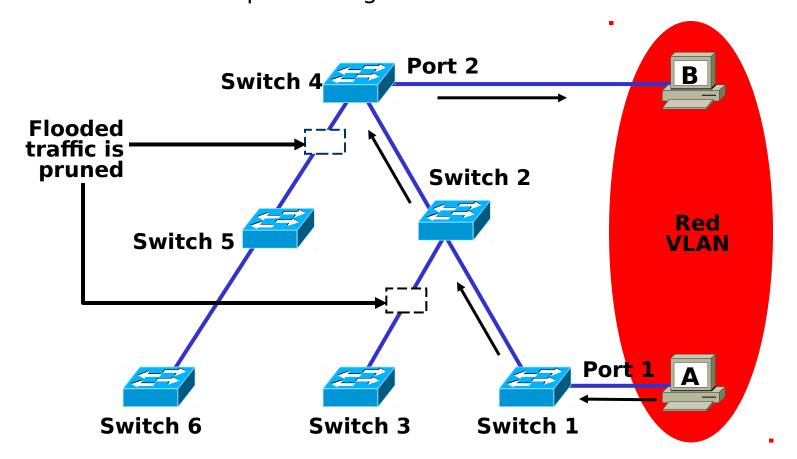


VTP Pruning



MSTP

Increases available bandwidth by reducing unnecessary flooded traffic Example: Station A sends a broadcast, broadcast is only flooded to switches that have ports assigned to the red VLAN



VLAN Guidelines





Maximum number of VLANs, switch-dependent Catalyst 6XXX, 55XX, 35XX, support 1005 VLANS Catalyst 1900 supports 64 VLANs

5 Factory Default VLANS VLAN 1 (Can not change VLAN 1 name) VLAN 1001-1005

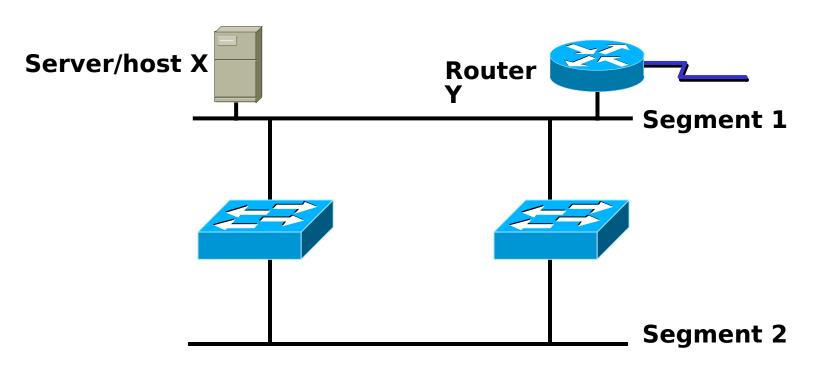
CDP and VTP advertisements are sent on VLAN1

Must be in VTP server or transparent mode to create, add, or delete VLANs

Redundant Topology





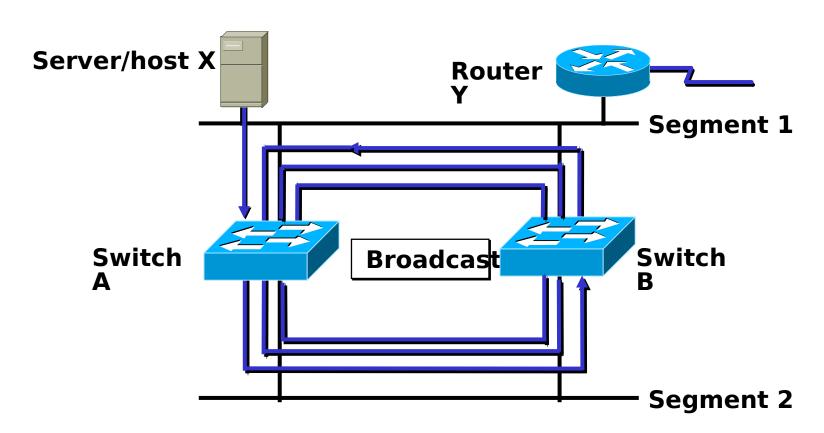


- Redundant topology eliminates single points of failure
- Redundant topology causes broadcast storms, multiple frame copies, and MAC address table instability problems

Broadcast Storms





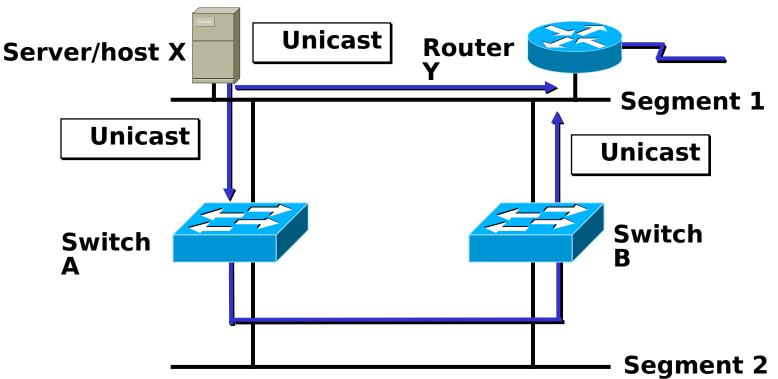


Switches continue to propagate broadcast traffic over and over

Multiple Frame Copies





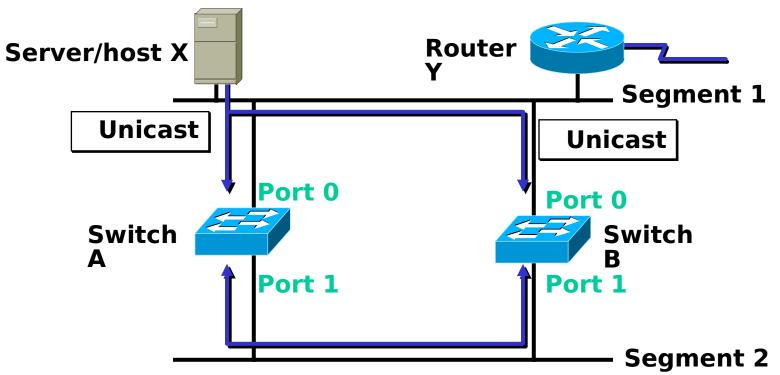


- Host X sends an unicast frame to Router Y
- Router Y MAC Address has not been learned by either Switch yet
- Router Y will receive two copies of the same frame



MAC Database Instability



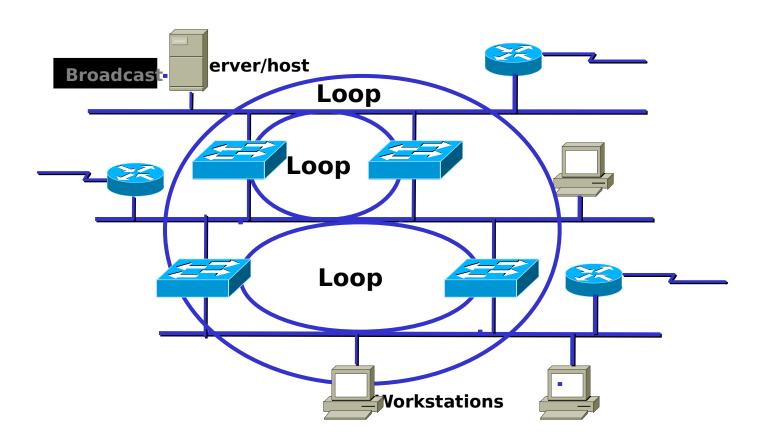


- Host X sends an unicast frame to Router Y
- Router Y MAC Address has not been learned by either Switch y
- Switch A and B learn Host X MAC address on port 0
- Frame to Router Y is flooded
- Switch A and B incorrectly learn Host X MAC address on port 1



Multiple Loop Problems



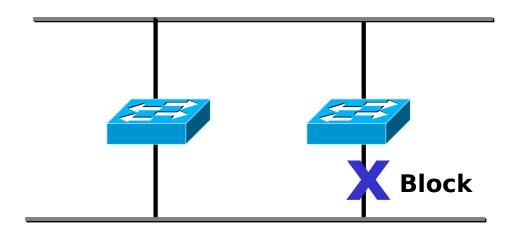


- Complex topology can cause multiple loops to occur
- Layer 2 has no mechanism to stop the loop

Solution: Spanning-Tree Protocol







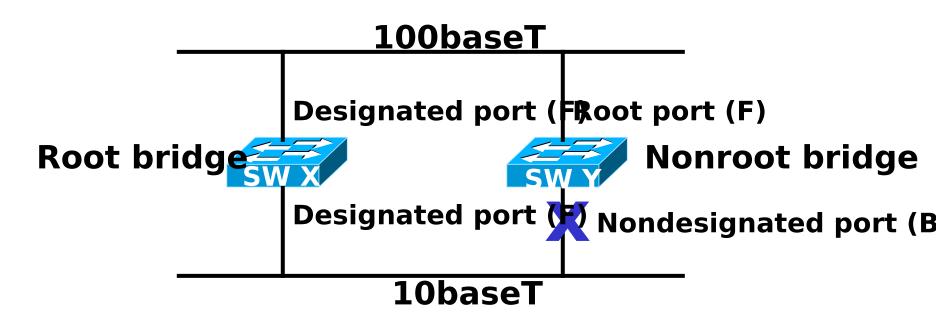
Provides a loop free redundant network topology placing certain ports in the blocking state



Spanning-Tree



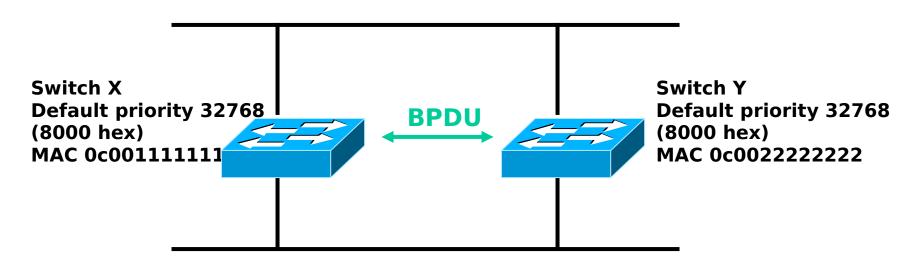
- Operations
- One root bridge per network
- One root port per nonroot bridge
- One designated port per segment



Root Bridge Selection



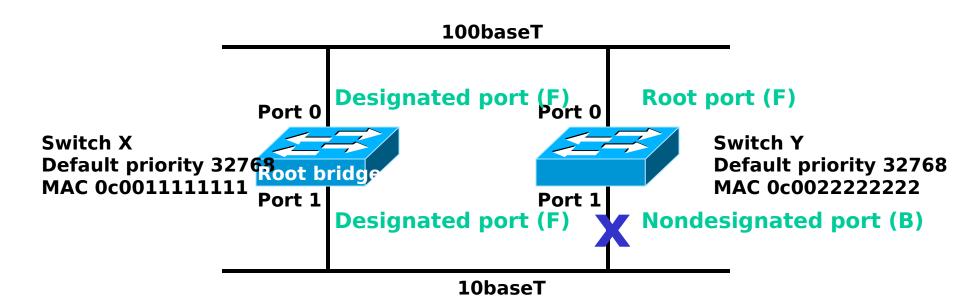








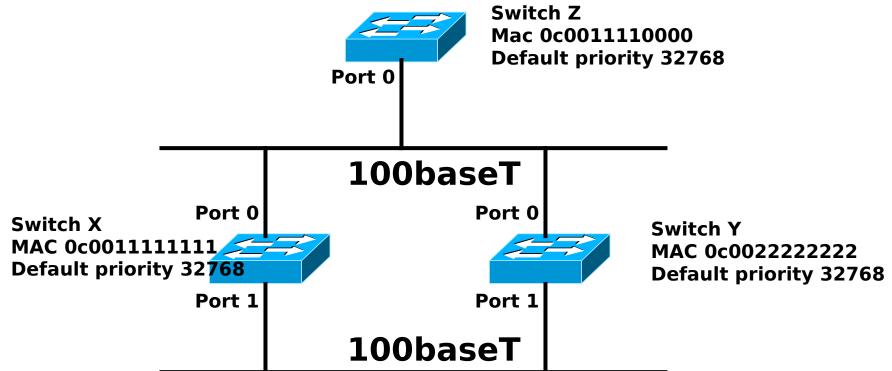




Spanning-Tree







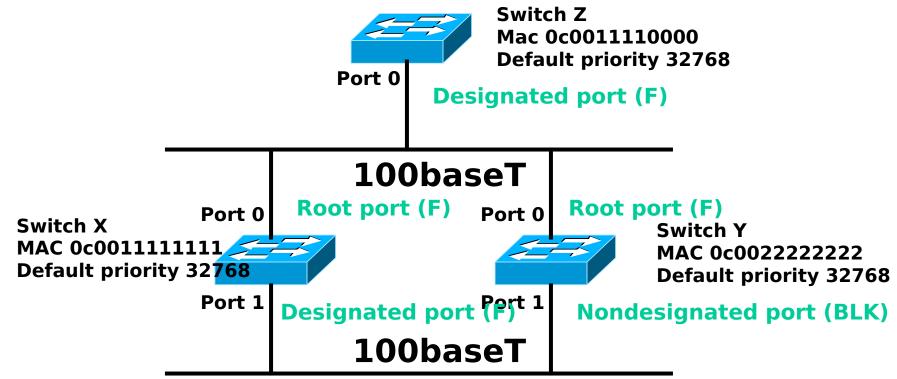
Can you figure out:

- What is the root bridge?
- What are the designated, nondesignated, and root parts?
- Which are the forwarding and blocking ports?

Spanning-Tree







Can you figure out:

- What is the root bridge?
- What are the designated, nondesignated, and root parts?
- Which are the forwarding and blocking ports?